Exam in Discrete Mathematics

First Year at the Faculty of Engineering and Science and the Technical Faculty of IT and Design

June 4th, 2019, 9.00-13.00

This exam consists of 11 numbered pages with 12 problems. All the problems are "multiple choice" problems. The answers must be given on these sheets.

It is allowed to use books, notes, photocopies etc. It is not allowed to use any electronic devices such as pocket calculators, mobile phones or computers.

The listed percentages specify by which weight the individual problems influence the total examination.

Remember to write your full name (including middle names) together with your student number below.

NAME:	
STUDENT NUMBER:	

There is only one correct answer to each question.

Problem 1 (8 %)

Let \mathcal{R} be a relation on the set of all integers and defined by $x \geq y^2$ and \mathcal{S} be a relation on the set of all integers and defined by $x = y^2$. What are the truth values of the following statements? 1. \mathcal{R} is transitive. True ☐ False 2. \mathcal{R} is antisymmetric. ☐ True ☐ False 3. \mathcal{R} is an equivalence relation. True ☐ False 4. $S \circ \mathcal{R} = \mathcal{R}$. ☐ True ☐ False **Problem 2** (8 %) Determine whether each of the following statements is true or false. 1. The set of integers with the absolute value less than 1000 is countable. ☐ TRUE ☐ FALSE 2. If A is an uncountable set and B is a countable set then A-B (the difference of A and B) must be uncountable. ☐ TRUE ☐ FALSE 3. It is true that $x^2 = y^2$ if and only if x = y or x = -y. ☐ TRUE ☐ FALSE

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4. If n is a positive integer, then n is odd if and only if 5n + 6 is odd.

☐ FALSE

☐ TRUE

1.	Which o	ne of the	following	integers x	satisfies 1	that	
				$x \equiv 4$	(mod 3) (mod 5) (mod 7)		
	<u> </u>	□ 34	7	□ 32	1 3	□ 24	14
2.	Find the			$ 27x \equiv 3 $ $ 363 $	•	,	☐ 141
			P	roblem 4	(8 %)		
Dete	rmine wh	ether eacl	n of the fo	ollowing th	eorems is	true or fa	alse.
1.	$(\mathbb{Z},=)$ is	a poset.					
	☐ TRU	E			☐ FALS	SE	

2. (\mathbb{Z}, \nmid) is a poset.

☐ TRUE

Problem 3 (8%)

☐ FALSE

			P	c meldor	(9	%)			
Ansv	ver the fol	llowing tr	ue/false p	roblems.	`	,			
1.	$2^x + 17$ i	is $O(3^x)$				True			False
2.	3x + 7 is	$\Omega(x)$				True			False
3.	$2x^2 + x$	-7 is $\Theta(a)$;2)			True			False
			Pı	roblem 6	(8	%)			
1.	What is	the value	of 4 ⁵³² m	od 11?					
	_ 2	3	<pre>1</pre>	<u> </u>		4	□ 8		7
2.	What is	50! mod	50?						
	□ 49	3	□ 7	□ 51		47	□ 46)
					/0	W.)			
			Pi	roblem 7	(8	%)			
Cons	ider the f	ollowing a	lgorithm:						
CC		no 1's yet		a_n : bit	strir	ng)			
re		= 1 then o		$\begin{array}{c} { m unt+1} \\ { m the \ numbe} \end{array}$	er of	f 1's}			
		•		nber of co			used by t	his a	algorithm.
] 0	$\left(\frac{\log n}{n}\right)$	$\bigcap O(\frac{1}{n})$)	\square $O(1)$			O(n)		$]$ $O(\sqrt{n})$

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Problem 8 (9 %)

1.	Find a recurrence relation for the number of bit strings of length n that contain three consecutive 0s.									
	$a_n = a_{n-1} + a_{n-2} + a_{n-3} + 2^{n-3}$ for all $n \ge 3$.									
	$a_n = a_{n-2} + a_{n-1} + 2^n \text{ for all } n \ge 2.$									
	$a_n = a_{n-1} + a_{n-2} + 2^{n-2}$ for all $n \ge 2$.									
	$a_n = a_{n-1} + a_{n-2} + a_{n-3}$ for all $n \ge 3$.									
2.	What are the initial conditions?									
	$\Box \ a_0 = a_1 = a_2 = 0.$									
	$a_0 = 1, a_1 = 2 \text{ og } a_2 = 0.$									
	$a_0 = 0, a_1 = 2 \text{ og } a_2 = 0.$									
	$a_0 = 1, a_1 = 2 \text{ og } a_2 = 2.$									
3.	How many bit strings of length five contain three consecutive 0s?									
	1 5 12 8 20 3 47									

Problem 9 (8 %)

- 1. What is the coefficient of $x^{101}y^{99}$ in $(2x-3y)^{200}$?
 - $\Box -2^{101}3^{99}C(200,99)$

 - $-2^{101}3^{99}P(200,99)$

 - C(200,99)
- 2. Suppose that k, n are integers with $1 \le k < n$. Which one is equal to

$$C(n-1, k-1) \cdot C(n, k+1) \cdot C(n+1, k)$$
?

- $C(n-1,k) \cdot C(n,k-1) \cdot C(n+1,k+1)$
- $\Box C(n,k) \cdot C(n-1,k-1) \cdot C(n,k+1)$
- $C(n-1, k-2) \cdot C(n+1, k+1)$

where P(n, k) is the k-permutation of a set with n elements and where C(n, k) is the k-combination of a set with n elements..

Problem 10 (8 %)

Theorem For every positive integer n, $\sum_{i=1}^n i = \frac{(n+\frac{1}{2})^2}{2}$. Basis Step: The formula is true for n=1. Inductive Step: Suppose that $\sum_{i=1}^n i = \frac{(n+\frac{1}{2})^2}{2}$. Then $\sum_{i=1}^{n+1} i = (\sum_{i=1}^n i) + (n+1)$. By the inductive hypothesis, $\sum_{i=1}^{n+1} i = \frac{(n+\frac{1}{2})^2}{2} + n + 1 = (n^2 + n + \frac{1}{4})/2 + n + 1 = (n^2 + 3n + \frac{9}{4})/2 = (n+\frac{3}{2})^2/2 = [(n+1) + \frac{1}{2}]^2/2$, completing the inductive step.

What is wrong with this "proof"?

	Basis	step	must	be	verified	for	n =	: 0.
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- ☐ The inductive step is wrong.
- Basis step is wrong.
- There is nothing wrong with this proof.

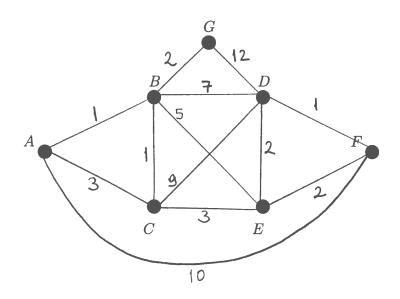


Figure 1: The graph G, considered in Problem 11.

Problem 11 (9%)

In this problem we use Dijkstra's algorithm (see Figure 2) on the graph G in Figure 1.

1.	What is the length	of the shortest pa	th from A to I	F?		
	□ 10	9	□ 8		7	
2.	In what order are Figure 2)?	vertices added to	the set S in I	Dijkstra's	algorithm	(see
		7				

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3.	What is	the	weight	of a	minimum	spanning	tree o	of the	graph	${\cal G}$ in	Figure	e 1.
	□ 8											
	7											
	□ 5											
	<u> </u>											

Problem 12 (9%)

Let K_n denote the complete, undirected graph with n nodes (i.e. every pair of the n-nodes contains an edge between them).

1.	When (for what values of n) does K_n h	ave an Euler path?
2.	When (for what values of n) does K_n h	ave a Hamilton path?
3.	Is it possible to have a graph that has a	a Hamilton path but no Euler path?
	☐ YES [NO

```
procedure Dijkstra(G: weighted connected simple graph, with
     all weights positive)
\{G \text{ has vertices } a = v_0, v_1, \dots, v_n = z \text{ and lengths } w(v_i, v_j)
     where w(v_i, v_j) = \infty if \{v_i, v_j\} is not an edge in G\}
for i := 1 to n
     L(v_i) := \infty
L(a) := 0
S := \emptyset
{the labels are now initialized so that the label of a is 0 and all
     other labels are \infty, and S is the empty set}
while z \notin S
     u := a vertex not in S with L(u) minimal
     S := S \cup \{u\}
     for all vertices v not in S
           if L(u) + w(u, v) < L(v) then L(v) := L(u) + w(u, v)
           {this adds a vertex to S with minimal label and updates the
           labels of vertices not in S}
return L(z) {L(z) = length of a shortest path from a to z}
```

Figure 2: