#### Exam in Discrete Mathematics

# First Year at The TEK-NAT Faculty June 10th, 2016, 9.00-13.00

This exam consists of 11 numbered pages with 16 problems. All the problems are "multiple choice" problems. The answers must be given on these sheets.

It is allowed to use books, notes, photocopies etc. It is not allowed to use any electronic devices such as pocket calculators, mobile phones or computers.

The listed percentages specify by which weight the individual problems influence the total examination.

Remember to write your full name (including middle names) together with your student number below.

NAME:	
CTUDENT NUMBED.	
STUDENT NUMBER:	

 $There\ is\ only\ one\ correct\ answer\ to\ each\ question.$ 

#### **Problem 1** (8 %)

relatio

Consider the following linear homogeneous recurrence relation								
$a_n = -a_{n-1} + 6a_{n-2}.$	$a_n = -a_{n-1} + 6a_{n-2}.$							
1. What is the degree of this recurrence relation?								
$\begin{array}{cccccccccccccccccccccccccccccccccccc$								
2. Which of the following is the solution of this recurrence relation ( $\alpha_1$ are constants)?	and $\alpha_2$							
Problem 2 (5 %)								
Consider the following algorithm:								
$\begin{aligned} & \mathbf{procedure} \ \textit{multiplications}(n: \ \textit{positive integer}) \\ & t := 1 \\ & \mathbf{for} \ i := 1 \ \mathbf{to} \ n \\ & \mathbf{for} \ j := 1 \ \mathbf{to} \ i \\ & t := 2 \cdot t \end{aligned}$								
The number of multiplications used by this algorithm is								
	;)							

Problem	3	(5)	%
	•	v	70

Which one of th	e following numbers i	is an inverse of 43 m	odulo 100?	
□ -3	□ 7	17	□ 27	
	$\operatorname{Prob}$	lem 4 (9 %)		
Consider the following	lowing algorithm:			
procedure $a$ $i := 0$ $x := n$ $s := 0$ while $i < n$ $i := i + 1$ $s := s + i - 1$ return $s$	dditions(n): positive i $+x$	nteger)		
in this algorithm in this algorithm is $i \le n$ or $i \le n$ or $i \le n$ or $i \le n$ or $i \le n$	orithm?	= n		юр

	Problem 5 (6 %)						
W	What is value of $(79 + 778 + 7777 \cdot 321) \mod 7$ ?						
	0	3	☐ 4	<u> </u>	□ 6		
		Problem	6 (8 %)				
Le	t $P(n)$ be the following st	tatement abo	ut the absolut	te value of rea	al numbers		
	$ x_1 + x_2 + \ldots + x_n  \le  x_1 $	$+ x_2 +\ldots$	$+ x_n $ , for all	real numbers	$x_1,\ldots,x_n$ .		
	om an exercise in the boo e want to prove by induct	-		` '			
W	hich one of the following	is a correct or	utline of the i	nductive step	?		
	Let $k \ge 1$ and assume that $P(k)$ is true. Then use $ x_1 + x_2 + x_3 \dots + x_{k+1}  \le  x_1 + x_2  +  x_3 + \dots + x_{k+1} $ and the induction hypothesis to prove $P(k+1)$ .						
	Let $k \geq 2$ and assume the Then use $ x_1 + x_2 + \dots $ induction hypothesis to p	$+ x_{k-1} + x_k  $	$\leq  x_1 + x_2 +$	$\dots + x_{k-1} $	$+  x_k $ and the		
	Let $k \geq 2$ and assume the Then use $ x_1 + x_2 + \dots $ induction hypothesis to p	$+  x_k + x_{k+1} $	$\leq  x_1 + x_2 +$	$\cdots + x_k  +  $	$ x_{k+1} $ and the		

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Then use  $|x_1 + x_2 + \ldots + x_{k+1}| \le |x_1| + |x_2 + \ldots + x_{k+1}|$  and the induction

Then use  $|x_1 + x_2 + \ldots + x_k + x_{k+1}| \le |x_1 + x_2 + \ldots + x_k| + |x_{k+1}|$  and the

 $\square$  Let  $k \geq 3$  and assume that P(k) is true.

 $\square$  Let  $k \geq 2$  and assume that P(k-1) is true.

induction hypothesis to prove P(k).

hypothesis to prove P(k+1).

#### Problem 7 (9%)

Consider the following two relations on the set  $A = \{1, 2, 3, 4, 5\}$ :

$$R = \{(1,1), (1,2), (2,3), (3,4), (4,5), (5,5)\}$$
 
$$S = \{(1,2), (1,3), (2,1), (3,1), (3,4), (3,5), (4,3), (4,4), (5,3)\}.$$

1.	. Answer the following true/false problems:						
	R is reflexive	e		☐ True		] False	
	R is antisym	nmetric		☐ True		] False	
	S is symmet	ric		☐ True		] False	
	S is transitiv	ve		☐ True		] False	
2.	Which one o relation $R \circ$		ng values of $x$	satisfies tha	at $(5,x)$ is	in the comp	osed
	<u> </u>	<u> </u>	<u> </u>		4	□ 5	
3.	Which one of relation $S \circ$		ng values of $x$	satisfies tha	at $(5,x)$ is	in the comp	osed
	<u> </u>	$\square$ 2	<u> </u>		4	□ 5	
4.	Let $R^*$ denotin $R^*$ ?	te the trans	itive closure o	of $S$ . How n	nany pairs	(a,b) are t	here
	□ 6	9	<b>1</b> 0	<u> </u>	<u> </u>	<u> </u>	

Problem 8 (	7	%)

A graph G with vertices  $v_1, v_2, \dots, v_7$  has adjacency matrix

$$\begin{bmatrix} 0 & 1 & 0 & 1 & 0 & 0 & 0 \\ 1 & 0 & 1 & 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 & 1 & 1 & 1 \\ 1 & 0 & 0 & 0 & 1 & 1 & 1 \\ 0 & 0 & 1 & 1 & 0 & 0 & 0 \\ 0 & 0 & 1 & 1 & 0 & 0 & 0 \\ 0 & 0 & 1 & 1 & 0 & 0 & 0 \end{bmatrix}$$

1. Answer the following true/false problems.								
G has an Euler circuit	☐ True	☐ False						
G has a Hamilton circuit	☐ True	☐ False						
G has a Hamilton path	☐ True	☐ False						
2. What is the length of the shortest sin	aple circuit in $G$ ?							
	5 [	6						
3. What is the number of edges of a spa	nning tree of $G$ ?							
	□ 5 □ 6 □	7 🗌 8 📗 9						
<b>Problem 9</b> (4 %)								
Which one of the following propositions is	equivalent to $\forall x (P($	$(x) \land \neg Q(x))$ ?						

## **Problem 10** (4 %)

Answer the following problems about nu	mbers of permutation	ns and combinations					
$\forall n \forall r (P(n,r) = P(n,n-r))$	☐ True	☐ False					
$\forall n \forall r (C(n,r) = C(n,n-r))$	☐ True	☐ False					
$\forall n \forall r (P(n,r) = C(n,r) \cdot r!)$	☐ True	☐ False					
$\forall n \forall r (C(n+1,r+1) = C(n,r))$	☐ True	☐ False					
Problem	<b>Problem 11</b> (4 %)						
Consider the following statements:							
Statement 1: If it is Sunday then it is	sunshine.						
Statement 2: If it is not Sunday then	it is not sunshine.						
Statement 3: If it is sunshine then it is	s Sunday.						
Statement 4: If it is not sunshine then	it is not Sunday.						
What is the contrapositive of Statement	1?						
☐ Statement 1 ☐ Statement 2	☐ Statement 3	Statement 4					

## **Problem 12** (6 %)

Let .	$A = \{\emptyset, \emptyset\}$	$2, \{1, 2\}$	and $B$	$= \{a, \{\emptyset\}\}$	$\{1, 2, 3\}$	be sets.			
1.	What i	is the car	rdinality	of $A \cap E$	3 ?				
	<b>0</b>	<u> </u>	$\square$ 2	<b>3</b>	□ 4	<u> </u>	☐ 6	□ 7	□ 8
2.	What i	is the car	rdinality	of $A \cup E$	3 ?				
	<b>0</b>	<u> </u>	$\square$ 2	□ 3	□ 4	□ 5	□ 6	□ 7	□ 8
3.	Which	one of t	he follow	ring is an	element	of $A \times I$	3 ?		
		1)		1, 2)		$\{1, 2\}$			
4.	Which	one of the	he follow	ring is an	element	of the p	ower set	$\mathcal{P}(A)$ ?	
	□ {{∅	}}	□ {	$1, 2$ }					
				Proble	e <b>m 13</b> (4	1 %)			
1.	What i	is the val	lue of $P($	(6,3) ?					
	☐ 6 ☐ 30 ☐ 120 ☐ 720								
2.		is the val	lue of $C($	(7,3)?					
	$ \begin{array}{c c}  & 15 \\  & 20 \\  & 35 \end{array} $								
	<u> </u>								

## **Problem 14** (6 %)

Let $f(x) = (2x+3)(x^3)$ Answer the following to			
1. $f(x)$ is $O(x^3)$		☐ True	☐ False
2. $f(x)$ is $O(x^4)$		☐ True	☐ False
3. $f(x)$ is $O(x^5)$		☐ True	☐ False
4. $f(x)$ is $\Omega(x^3)$		☐ True	☐ False
5. $f(x)$ is $\Omega(x^4)$		☐ True	☐ False
6. $f(x)$ is $\Omega(x^5)$		☐ True	☐ False
7. $f(x)$ is $\Theta(x^3)$		☐ True	☐ False
8. $f(x)$ is $\Theta(x^4)$		☐ True	☐ False
9. $f(x)$ is $\Theta(x^5)$		☐ True	☐ False
	Problem 15	(10 %)	
In this problem we use	Dijkstra's algorithm (	see Figure 2) on the	graph in Figure 1.
1. What is the leng algorithm)?	th of the shortest p	ath from $a$ to $z$ (for	ound by Dijkstra's
☐ 6	7 🔲 8	<u> </u>	☐ 14
2. In what order are $a, v_1, v_2, v_5, v_8$ $a, v_1, v_4, v_5, v_8$ $a, v_1, v_3, v_2, v_4, a, v_1, v_2, v_3, v_4, a, v_1, v_2, v_3, v_4, a$		ne set $S$ ?	

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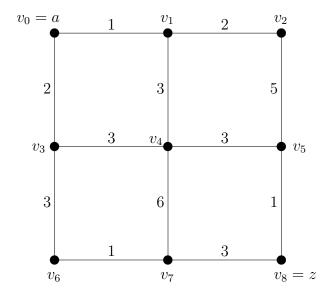


Figure 1: The graph G, considered in problems 15 and 16.

#### **Problem 16** (5 %)

What is	the weight	of a mini	mum span	ning tree o	of the grap	h in Figure	e 1.
9	11	□ 12	13	☐ 14	□ 15	☐ 16	18

```
procedure Dijkstra(G: weighted connected simple graph, with
     all weights positive)
{G has vertices a = v_0, v_1, \dots, v_n = z and lengths w(v_i, v_j)
     where w(v_i, v_j) = \infty if \{v_i, v_j\} is not an edge in G\}
for i := 1 to n
     L(v_i) := \infty
L(a) := 0
S := \emptyset
{the labels are now initialized so that the label of a is 0 and all
     other labels are \infty, and S is the empty set}
while z \notin S
     u := a vertex not in S with L(u) minimal
     S := S \cup \{u\}
     for all vertices v not in S
           if L(u) + w(u, v) < L(v) then L(v) := L(u) + w(u, v)
           {this adds a vertex to S with minimal label and updates the
           labels of vertices not in S}
return L(z) {L(z) = length of a shortest path from a to z}
```

Figure 2: