Exam in Discrete Mathematics

First Year at the Faculty of Engineering and Science and the Technical Faculty of IT and Design

June 4th, 2018, 9.00-13.00

This exam consists of 11 numbered pages with 14 problems. All the problems are "multiple choice" problems. The answers must be given on these sheets.

It is allowed to use books, notes, photocopies etc. It is not allowed to use any electronic devices such as pocket calculators, mobile phones or computers.

The listed percentages specify by which weight the individual problems influence the total examination.

Remember to write your full name (including middle names) together with your student number below.

NAME:	
STUDENT NUMBER:	

 $There\ is\ only\ one\ correct\ answer\ to\ each\ question.$

Problem 1 (6 %)

1100	(0 /0)	
$P(x,y)$ means " $x \leq y$ ", where the cointegers. What are the truth values of		
1. $\forall n P(0,n)$	☐ True	☐ False
$2. \ \forall x \exists y P(x,y)$	☐ True	☐ False
3. $\exists y \forall x P(x,y)$	☐ True	☐ False
Prob	olem 2 (8 %)	
Determine whether each of the follow	ving statements is tru	e or false.
1. If A and B are sets and $A \subseteq B$	$3 \text{ then } A \le B .$	
☐ TRUE	☐ FALSE	
2. If A and B are sets, A is uncountable.	intable, and $A \subseteq B$ the	nen B is countable.
☐ TRUE	☐ FALSE	
3. If A and B are sets with $ A =$	$ B $ then $ \mathcal{P}(A) = \mathcal{P}(A) $	$\mathcal{C}(B)$.
☐ TRUE	☐ FALSE	
4. If A is an infinite set, then it c	ontains a countably in	nfinite subset.
☐ TRUE	☐ FALSE	

Problem 3 (8%)

Determine whether each of the following theorems is true or false. Assume that $a,b,c,m\in\mathbb{Z}$ with m>1

1. If $a \equiv b \pmod{m}$ and $a \equiv c \pmod{m}$ then $a \equiv b + c \pmod{m}$							
	☐ TRUE	☐ FALSE					
2.	If $a \equiv b \pmod{m}$ then $2a \equiv 2b \pmod{m}$	1 m).					
	☐ TRUE	☐ FALSE					
3.	If $a \equiv b \pmod{m}$ then $a \equiv b \pmod{2}$	(m).					
	☐ TRUE	☐ FALSE					
4.	If $a \equiv b \pmod{m^2}$ then $a \equiv b \pmod{m^2}$	m).					
	☐ TRUE	☐ FALSE					

Problem 4 (9 %) $\log x$), for x > 0. A

Let $f(x) = 0$ problems.	$(x! + 2^x)(x$	$c^3 + 2\log x$	(x) , for $x \geq x$	> 0.	Ansv	wer the fe	ollowing	true/false	
1. $f(x)$ is	1. $f(x)$ is $O(2^x x^3)$						☐ False		
2. $f(x)$ is	$O(x! \ x^3)$			Γ	True		☐ Fals	e	
3. $f(x)$ is	$O(\log x!)$			Γ	True		☐ Fals	e	
4. $f(x)$ is	$\Omega(2^x \ x^3)$			Γ	True		☐ Fals	e	
5. $f(x)$ is	$\Omega(x! \ x^3)$			☐ True ☐ Fa			☐ Fals	e	
6. $f(x)$ is	$\Omega(\log x!)$			Γ	True		☐ Fals	e	
7. $f(x)$ is	$\Theta(2^x \ x^3)$			☐ True			☐ Fals	e	
8. $f(x)$ is		☐ True			☐ Fals	e			
9. $f(x)$ is	$\Theta(\log x!)$			Γ	True		☐ Fals	e	
		Pı	oblem 5	(8 %	(o)				
1. What is	s the value	of 9 ⁴⁵ m c	od 23?						
□ 7	<u> </u>	<u> </u>	9	4		□ 3	□ 6		
2. What is	$s\ gcd(2^{89}, 2^{89})$	(2^{346}) ?							
2^3	$\square 2^4$	$ 2^{89}$	3^9	4		$ 2^{346}$	☐ 6		

Problem 6 (5 %)

Consider the following recursive algorithm: **procedure** power(n: nonnegative integer)if n = 0 then power(n) := 3else power $(n) := power(n-1) \cdot power(n-1)$ **return** power (n)Which one can be calculated using this recursive algorithm? 3^{2^n} 3^{2n} $\prod 3^n \cdot 2^{n-1}$ $\prod 2^n \cdot 3$ $\prod 3n$ **Problem 7** (6 %) We consider the following moves of a particle in the xy plane $R:(x,y) \rightarrow (x+1,y)$ (one unit right) $U:(x,y) \rightarrow (x,y+1)$ (one unit up)? In how many ways can the particle move from the origin to the point (8,5)? | P(8,5) | $\Box \ C(8,5)$ \Box C(13,8) $\prod P(13,5)$ $\prod P(13,8)$

Problem 8 (6 %)

What is the binomial expansion of $(x-\frac{3}{x})^5$?

- $x^5 + 15x^3 + 90x + \frac{270}{x} + \frac{405}{x^3} + \frac{243}{x^5}$
- $x^5 5x^3 + 15x \frac{45}{x} + \frac{45}{x^3} \frac{24}{x^5}$

Problem 9 (8 %)

- 1. What is a recurrence relation for the number of bit strings of length n that do not contain 3 consecutive 0's?
 - $a_n = 2a_{n-1} + a_{n-2} \text{ for } n \ge 3$

 - $\begin{array}{l} \square & a_n = 5a_{n-1} + a_{n-2} \text{ for } n \ge 3 \\ \square & a_n = 5a_{n-1} a_{n-3} \text{ for } n \ge 4 \\ \square & a_n = a_{n-1} + a_{n-2} + a_{n-3} \text{ for } n \ge 4 \\ \square & a_n = 2a_{n-1} + a_{n-2} + a_{n-3} \text{ for } n \ge 4 \\ \square & a_n = 2a_{n-1} + a_{n-2} \text{ for } n \ge 3 \end{array}$
- 2. What are the initial conditions?
 - $a_1 = 2, a_2 = 4, a_3 = 7$
 - $\boxed{} a_1 = 1, a_2 = 3, a_3 = 5$
 - $\begin{array}{c}
 a_1 = 2, a_2 = 4, a_3 = 6 \\
 a_1 = 2, a_2 = 3, a_3 = 7
 \end{array}$

Problem 10 (7 %)

Let P(n) be the following statement

$$\sum_{k=0}^{n} 3^k = \frac{3^{n+1} - 1}{2}$$

We want to prove by induction that P(n) is true for all $n \geq 0$.

1.	What	is	the	correct	basis	step	of	the	induction	proof?
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- \square Prove that P(1) is true
- \square Prove that P(2) is true
- \square Prove that P(0) is true
- \square Prove that P(n) is true, for all $n \ge 1$

2. Which one of the following is a correct outline of the inductive step?

- Let $i \ge -1$ and assume that P(i) is true. By the induction hypothesis $\sum_{k=0}^{i} 3^k = \frac{3^{i+1}-1}{2}$. Use this to prove P(i+2).
- Let $i \geq 0$ and assume that P(i) is true. By the induction hypothesis $\sum_{k=0}^{i} 3^k = \frac{3^{i+1}-1}{2}$. Use this to prove P(i+1).
- Let i=0 and assume that P(i) is true. By the induction hypothesis $\sum_{k=0}^{i} 3^k = \frac{3^{i+1}-1}{2}$. Use this to prove P(i).
- Let $i \geq 0$ and assume that P(i) is true. By the induction hypothesis $\sum_{k=0}^{i} 3^k = \frac{3^{i+1}-1}{2}$. Use this to prove P(i+2).
- Let $i \geq 2$ and assume that P(i) is true. By the induction hypothesis $\sum_{k=0}^{i} 3^k = \frac{3^{i+1}-1}{2}$. Use this to prove P(i+1).

Problem 11 (9 %)

Let \mathcal{R} be the relation on the set $A = \{1, 2, 3, 4\}$ defined by

$$\mathcal{R} = \{(i, j) \mid 2 \text{ divides } (i - j)\}.$$

1. What is the matrix that represents the relation \mathcal{R} ?

$$\square \begin{pmatrix} 1 & 0 & 1 & 1 \\ 0 & 1 & 0 & 0 \\ 1 & 0 & 1 & 1 \\ 1 & 0 & 1 & 1 \end{pmatrix} \qquad \square \begin{pmatrix} 1 & 1 & 1 & 1 \\ 0 & 1 & 0 & 0 \\ 1 & 0 & 1 & 1 \\ 1 & 0 & 1 & 1 \end{pmatrix} \qquad \square \begin{pmatrix} 1 & 0 & 1 & 0 \\ 0 & 1 & 0 & 1 \\ 1 & 0 & 1 & 0 \\ 0 & 1 & 0 & 1 \end{pmatrix}$$

2. Answer the following true/false problems.

\mathcal{R} is reflexive	☐ True	☐ False
\mathcal{R} is symmetric	☐ True	☐ False
\mathcal{R} is antisymmetric	☐ True	☐ False
\mathcal{R} is transitive	☐ True	☐ False
\mathcal{R} is an equivalence relation	□ True	☐ False

3. What is the matrix that represents the relation \mathbb{R}^2 ?

$$\Box \begin{pmatrix} 1 & 0 & 1 & 0 \\ 0 & 1 & 0 & 1 \\ 1 & 0 & 1 & 0 \\ 0 & 1 & 0 & 1 \end{pmatrix} \qquad \Box \begin{pmatrix} 1 & 1 & 1 & 1 \\ 0 & 1 & 0 & 0 \\ 1 & 0 & 1 & 1 \\ 1 & 0 & 1 & 1 \end{pmatrix} \qquad \Box \begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \end{pmatrix}$$

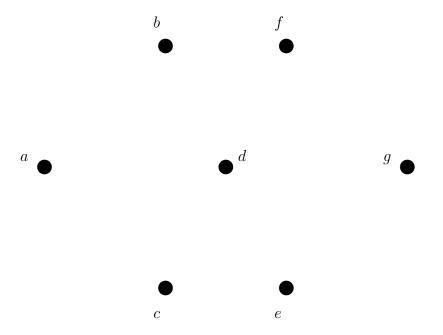


Figure 1: The graph G, considered in Problems 12, 13 and 14.

Problem 12 (8 %)

In this problem we use Dijkstra's algorithm (see Figure 2 on page 11) on the graph G in Figure 1.

1.	hat is gorithn		th of the	shortest 1	path from	a to g (f	found by	Dijkstra's
	7	□ 8	9	□ 10	11	□ 12	□ 13	☐ 14
2.	a, c, d a, b, c, a, b, c, a, e, b, a, c, b, a, e, f	, e, g , f, d, e, g , d, e, f, g	vertices a	added to t	he set S	?		

Page 9 of 11

Problem 13 (5%)

What is the weight of a minimum spanning tree of the graph G in Figure 1.								
☐ 14	<u> </u>	□ 16	□ 17	□ 12	19	10	□ 13	
			Problem	14 (7 %))			
In this problem G is the graph in Figure 1. (The edge weights of G are not considered in this problem.)								
1. A	nswer the f	ollowing tr	rue/false pr	oblems.				
G	has an Eu	ler circuit		☐ Tr	rue	☐ Fals	se	
G	has an Eu	ler path		☐ Tr	rue	☐ Fals	se	
G has a Hamilton circuit $\ \square$ True $\ \square$ False							se	
G	has a Ham	nilton path		☐ Tr	rue	☐ Fals	se	
2. W	hat is the	length of a	shortest si	imple circu	uit of G ?			
	1 🔲	2 🔲 3	\Box 4	□ 5	□ 6	□ 7	□ 8	

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procedure Dijkstra(G: weighted connected simple graph, with
     all weights positive)
{G has vertices a = v_0, v_1, \dots, v_n = z and lengths w(v_i, v_j)
     where w(v_i, v_j) = \infty if \{v_i, v_j\} is not an edge in G\}
for i := 1 to n
     L(v_i) := \infty
L(a) := 0
S := \emptyset
{the labels are now initialized so that the label of a is 0 and all
     other labels are \infty, and S is the empty set}
while z \notin S
     u := a vertex not in S with L(u) minimal
     S := S \cup \{u\}
     for all vertices v not in S
           if L(u) + w(u, v) < L(v) then L(v) := L(u) + w(u, v)
           {this adds a vertex to S with minimal label and updates the
           labels of vertices not in S}
return L(z) {L(z) = length of a shortest path from a to z}
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Figure 2: